Chapter 1. Logic and Proofs

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Chapter 1. Logic and Proofs

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Definition

A **proposition** is a sentence that has exactly one truth value. It is either true, which we denote by T, or false, which we denote by F.

Example

$$7^2 > 60$$
 (F), $\pi > 3$ (T), Earth is the closest planet to the sun (F).

Example

The statement "the north Pacific right whale (露 脊 鯨) will be extinct species before the year 2525" has one truth value but it takes time to determine the truth value.

Example

That "Euclid was left-handed" is a statement that has one truth value but may never be known.

Definition

A **negation** of a proposition P, denoted by $\sim P$, is the proposition "not P". The proposition $\sim P$ is $\begin{array}{c} \mathsf{true} \\ \mathsf{false} \end{array} \text{ exactly when } P \text{ is } \begin{array}{c} \mathsf{false} \\ \mathsf{true} \end{array}.$

Definition

 $\begin{array}{c} \mbox{Given propositions P and Q, the} & \begin{array}{c} \mbox{conjunction} \\ \mbox{disjunction} \end{array} \mbox{of P and Q, denoted} \\ \mbox{by } \begin{array}{c} P \land Q \\ P \lor Q \end{array}, \mbox{ is the proposition "P} & \begin{array}{c} \mbox{and} \\ \mbox{and} \\ \mbox{or} \end{array} & \begin{array}{c} P \land Q \\ \mbox{or} \end{array} & \begin{array}{c} P \land Q \\ \mbox{P} \lor Q \end{array} \mbox{ is true exactly} \\ \mbox{both P and Q are true} \\ \mbox{at least one of P or Q is true} \end{array}$

Example

Now we analyze the sentence "either 7 is prime and 9 is even, or else 11 is not less than 3". Let P denote the sentence "7 is a prime", Q denote the sentence "9 is even", and R denote the sentence "11 is less than 3". Then the original sentence can be symbolized by $(P \wedge Q) \lor (\sim R)$, and the table of truth value for this sentence is

Р	Q	R	$P \wedge Q$	$\sim R$	$(P \land Q) \lor (\sim R)$
Т	Т	Т	Т	F	Т
T	Т	F	Т	Т	Т
Т	F	Т	F	F	F
F	Т	Т	F	F	F
Т	F	F	F	Т	Т
F	Т	F	F	Т	Т
F	F	Т	F	F	F
F	F	F	F	Т	Т

Since P is true and Q, R are false, the sentence $(P \land Q) \lor (\sim R)$ is true.

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Example

The logic symbol $\sim (P \lor \sim P) \lor (Q \land \sim Q)$ is a contradiction.

Definition

Two propositional forms are said to be *equivalent* if they have the same truth value.

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Theorem

For propositions P, Q, R, we have the following: (a) $P \Leftrightarrow \sim (\sim P)$. (Double Negation Law) $\begin{array}{c} (b) \ P \lor Q \Leftrightarrow Q \lor P \\ (c) \ P \land Q \Leftrightarrow Q \land P \end{array} \right\} \quad (\textbf{Commutative Laws})$ (d) $P \lor (Q \lor R) \Leftrightarrow (P \lor Q) \lor R$ (e) $P \land (Q \land R) \Leftrightarrow (P \land Q) \land R$ (Associative Laws) $\begin{array}{l} (f) \ P \land (Q \lor R) \Leftrightarrow (P \land Q) \lor (P \land R) \\ (g) \ P \lor (Q \land R) \Leftrightarrow (P \lor Q) \land (P \lor R) \end{array} \right\}$ (Distributive Laws) $\begin{array}{c} (h) \sim (P \land Q) \Leftrightarrow (\sim P) \lor (\sim Q) \\ (i) \sim (P \lor Q) \Leftrightarrow (\sim P) \land (\sim Q) \end{array} \right\}$ (De Morgan's Laws)

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Proof.

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We prove (g) for example, and the other cases can be shown in a similar fashion. Using the truth table,

Р	Q	R	$Q \land R$	P∨(Q∧R)	$P \lor Q$	P∨R	$(P \lor Q) \land (P \lor R)$
Т	Т	Т	Т	Т	Т	Т	Т
Т	Т	F	F	Т	Т	Т	Т
Т	F	Т	F	Т	Т	Т	Т
F	Т	Т	Т	Т	Т	Т	Т
Т	F	F	F	Т	Т	Т	Т
F	Т	F	F	F	Т	F	F
F	F	Т	F	F	F	Т	F
F	F	F	F	F	F	F	F
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Definition

A *denial* of a proposition is any proposition equivalent to $\sim P$.

- Rules for \sim , \wedge and $\vee\colon$
 - ${\rm O}~\sim$ is always applied to the smallest proposition following it.
 - $\mathbf{2}$ \wedge connects the smallest propositions surrounding it.
 - O v connects the smallest propositions surrounding it.

Example

Under the convention above, we have

$$2 P \lor Q \lor R \Leftrightarrow (P \lor Q) \lor R \Leftrightarrow P \lor (Q \lor R).$$

Definition

For propositions P and Q, the *conditional sentence* $P \Rightarrow Q$ is the proposition "if P, then Q". Proposition P is called the *antecedent* and Q is the *consequence*. The sentence $P \Rightarrow Q$ is true if and only if P is false or Q is true.

Remark:

In a conditional sentence, P and Q might not have connections. The truth value of the sentence "P \Rightarrow Q" only depends on the truth value of P and Q.

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Example

We would like to determine the truth value of the sentence "if x > 8, then x > 5". Let P denote the sentence "x > 8" and Q the sentence "x > 5".

- If P, Q are both true statements, then x > 8 which is (exactly the same as P thus) true.
- ② If P is false while Q is true, then 5 < x ≤ 8 which is (exactly the same as $\sim P \land Q$ thus) true.
- If P, Q are both false statements, then x ≤ 5 which is (exactly the same as ~Q thus) true.
- ${\ensuremath{\textcircled{}}}$ It is not possible to have P true but Q false.

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• How to read $\mathrm{P} \Rightarrow \mathrm{Q}$ in English?

- 1. If P, then Q. \quad 2. P is sufficient for Q. \quad 3. P only if Q.
- 4. Q whenever P.~ 5. Q is necessary for P.~ 6. Q, if/when P.

Definition

Let \boldsymbol{P} and \boldsymbol{Q} be propositions.

- The *converse* of $P \Rightarrow Q$ is $Q \Rightarrow P$.
- 2 The *contrapositive* of $P \Rightarrow Q$ is $\sim Q \Rightarrow \sim P$.

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Example

We would like to determine the truth value, as well as the converse and the contrapositive, of the sentence "if π is an integer, then 14 is even".

- Since that π is an integer is false, the implication "if π is an integer, then 14 is even" is true.
- 2 The converse of the sentence is "if 14 is even, then π is an integer" which is a false statement.
- The contrapositive of the sentence is "if 14 is not even, then π is not an integer" which is a true statement since the antecedent "14 is not even" is false.

By this example, we know that a sentence and its converse cannot be equivalent.

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Theorem

For propositions P and Q, the sentence $P\Rightarrow Q$ is equivalent to its contrapositive $\sim Q\Rightarrow \sim P.$

Proof.

Using the truth table



we conclude that the truth value of $P\Rightarrow Q$ and $\sim Q\Rightarrow \sim P$ are the same; thus they are equivalent sentences. $\hfill\square$

Definition

For propositions P and Q, the **bi-conditional sentence** $P \Leftrightarrow Q$ is the proposition "P if and only if Q". The sentence $P \Leftrightarrow Q$ is true exactly when P and Q have the same truth values. In other words, $P \Leftrightarrow Q$ is true if and only if P is equivalent to Q.

Remark: The notation \Leftrightarrow is a combination of \Rightarrow and its converse \Leftarrow , so the notation seems to suggest that $(P \Leftrightarrow Q)$ is equivalent to $(P \Rightarrow Q) \land (Q \Rightarrow P)$. This is in fact true since



Example

- The proposition " $2^3 = 8$ if and only if 49 is a perfect square" is true because both components are true.
- 2 The proposition " $\pi = \frac{22}{7}$ if and only if $\sqrt{2}$ is a rational number" is also true (since both components are false).
- The proposition "6 + 1 = 7 if and only if Argentina is north of the equator" is false because the truth values of the components differ.

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Remark:

Definitions may be stated with the "if and only if" wording, but it is also common practice to state a formal definition using the word "if". For example, we could say that "a function f is continuous at a number c if \cdots " leaving the "only if" part understood.

Example

A teacher says "If you score 74% or higher on the next test, you will pass the exam". Even though this is a conditional sentence, everyone will interpret the meaning as a biconditional (since the teacher tries to "define" how you can pass the exam).

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Theorem

For propositions P, Q and R, we have the following:

(a)
$$(P \Rightarrow Q) \Leftrightarrow (\sim P \lor Q)$$
.
(b) $(P \Leftrightarrow Q) \Leftrightarrow (P \Rightarrow Q) \land (Q \Rightarrow P)$.
(c) $\sim (P \Rightarrow Q) \Leftrightarrow (P \land \sim Q)$.
(d) $\sim (P \land Q) \Leftrightarrow (P \Rightarrow \sim Q)$.
(e) $\sim (P \land Q) \Leftrightarrow (Q \Rightarrow \sim P)$.
(f) $P \Rightarrow (Q \Rightarrow R) \Leftrightarrow (P \land Q) \Rightarrow R$.
(g) $P \Rightarrow (Q \land R) \Leftrightarrow (P \Rightarrow Q) \land (P \Rightarrow R)$.
(h) $(P \lor Q) \Rightarrow R \Leftrightarrow (P \Rightarrow R) \land (Q \Rightarrow R)$.

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- How to read $P \Leftrightarrow Q$ in English?
 - 1. P if and only if Q. 2. P if, but only if, Q.
 - 3. P implies Q, and conversely. 4. P is equivalent to Q.
 - 5. P is necessary and sufficient for Q.
- Rules for \sim , \wedge , \vee , \Rightarrow and \Leftrightarrow : These connectives are always applied in the order listed.

Example

- $\ \, \textbf{O} \ \, P \lor \sim Q \Leftrightarrow R \Rightarrow S \text{ is an abbr. for } \left[P \lor (\sim Q) \right] \Leftrightarrow (R \Rightarrow S).$
- $\textcircled{O} P \Rightarrow Q \Rightarrow R \text{ is an abbr. for } (P \Rightarrow Q) \Rightarrow R.$

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Definition

An **open sentence** is a sentence that contains variables. When P is an open sentence with a variable x (or variables x_1, \dots, x_n), the sentence is symbolized by P(x) (or $P(x_1, \dots, x_n)$). The **truth set** of an open sentence is the collection of variables (from a certain universe) that may be substituted to make the open sentence a true proposition. (使得 P(x) 為真的所有 x 形成 the truth set of P(x))

Remark:

In general, **an open sentence is not a proposition**. It can be true or false depending on the value of variables.

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Example

Let P(x) be the open sentence "x is a prime number between 5060 and 5090". In this open sentence, the universe is usually chosen to be \mathbb{N} , the natural number system, and the truth set of P(x) is $\{5077, 5081, 5087\}$.

Remark:

The truth set of an open sentence P(x) depends on the universe where x belongs to. For example, suppose that P(x) is the open sentence " $x^2 + 1 = 0$ ". If the universe is \mathbb{R} , then P(x) is false for all x (in the universe). On the other hand, if the universe is \mathbb{C} , the complex plane, then P(x) is true when $x = \pm i$ (which also implies that the truth set of P(x) is $\{i, -i\}$).

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Definition

With a universe X specified, two open sentences P(x) and Q(x) are equivalent if they have the same truth set of all $x \in X$.

Example

The two sentences "3x + 2 = 20" and "2x - 7 = 5" are equivalent open sentences in any of the number system, such as \mathbb{N} , \mathbb{Z} , \mathbb{Q} , \mathbb{R} and \mathbb{C} .

Example

The two sentences " $x^2 - 1 > 0$ " and " $(x < -1) \lor (x > 1)$ " are equivalent open sentences in \mathbb{R} .

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Given an open sentence P(x), the first question that we should ask ourself is "whether the truth set of P(x) is empty or not".

Definition

The symbol \exists is called the *existential quantifier*. For an open sentence P(x), the sentence $(\exists x)P(x)$ is read "there exists x such that P(x)" or "for some x, P(x)". The sentence $(\exists x)P(x)$ is true if the truth set of P(x) is non-empty.

Remark:

An open sentence P(x) does **not** have a truth value, but the quantified sentence $(\exists x)P(x)$ does.

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Example

The quantified sentence $(\exists x)(x^7 - 12x^3 + 16x - 3 = 0)$ is true in the universe of real numbers.

Example (Fermat number)

The quantified sentence $(\exists n)(2^{2^n} + 1 \text{ is a prime number})$ is true in the universe of natural numbers.

Example (Fermat's last theorem)

The quantified sentence

$$(\exists x, y, z, n)(x^n + y^n = z^n \land n \ge 3)$$

is true in the universe of integers, but is false in the universe of natural numbers.

Definition

The symbol \forall is called the *universal quantifier*. For an open sentence P(x), the sentence $(\forall x)P(x)$ is read "for all x, P(x)", "for every x, P(x)" or "for every given x (in the universe), P(x)". The sentence $(\forall x)P(x)$ is true if the truth set of P(x) is the entire universe.

Example

The quantified sentence $(\forall \, n)(2^{2^n}+1 \text{ is a prime number})$ is false in the universe of natural numbers since

$$2^{2^6} + 1 = 641 \times 6700417 \,.$$

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In general, statements of the form "every element of the set A has the property P" and "some element of the set A has property P" may be symbolized as $(\forall x \in A)P(x)$ and $(\exists x \in A)P(x)$, respective. Moreover,

● "All P(x) are Q(x)" (<u>所有滿足 P 的 x 都滿足 Q</u> or <u>只要滿</u> <u>足 P 的 x 就滿足 Q</u>) should be symbolized as "($\forall x$)(P(x) ⇒ Q(x))".

(See the next slide for the explanation!)

 [●] "Some P(x) are Q(x)" (<u>有些满足 P 的 x 也满足 Q</u> or <u>有些</u> <u>x 同時满足 P 和 Q</u>) should be symbolized as

" $(\exists x) (P(x) \land Q(x))$ ".

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• Explanation of 1: Suppose that the truth set of P(x) is A and the truth set of Q(x) is B. Then "All P(x) are Q(x)" implies that $A \subseteq B$; that is, if x in A, then x in B. Therefore, by reading the truth table

$x \in A$	$x \in B$	P(x)	Q(x)	$P(x) \Rightarrow Q(x)$
Т	Т	Т	Т	Т
Т	F	Т	F	F
F	Т	F	Т	Т
F	F	F	F	Т

we find that the truth set of the open sentence $P(x) \Rightarrow Q(x)$ is the whole universe since the second case $(x \in A) \land \sim (x \in B)$ cannot happen.

A (1) > A (2) > A (2) >

Example

• The sentence "for every odd prime x less than 10, $x^2 + 4$ is prime" can be symbolized as

 $(\forall x) [(x \text{ is odd}) \land (x \text{ is prime}) \land (x < 10) \Rightarrow (x^2 + 4 \text{ is prime})].$

The sentence "for every rational number there is a larger integer" can be symbolized as

$$(\forall x \in \mathbb{Q}) [(\exists z \in \mathbb{Z})(z > x)].$$

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Example

• The sentence "some functions defined at 0 are not continuous at 0" can be symbolized as

 $(\exists f)[(f \text{ is defined at } 0) \land (f \text{ is not continuous at } 0)].$

The sentence "some integers are even and some integers are odd" can be symbolized as

 $(\exists x)(x \text{ is even}) \land (\exists y)(y \text{ is odd}).$

 The sentence "some real numbers have a multiplicative inverse" (有些實數有乘法反元素) can be symbolized as

$$(\exists x \in \mathbb{R}) [(\exists y \in \mathbb{R})(xy = 1)].$$

To symbolized the sentence "any real numbers have an additive inverse" (任何實數都有加法反元素), it is required that we combine the use of the universal quantifier and the existential quantifier:

$$(\forall x \in \mathbb{R}) [(\exists y \in \mathbb{R})(x + y = 0)].$$

This is in fact quite common in mathematical statement. Another example is the sentence "some real number does not have a multiplicative inverse" (有些實數沒有乘法反元素) which can be symbolized by

$$(\exists x \in \mathbb{R}) \sim [(\exists y \in \mathbb{R})(xy = 1)]$$

or simply

$$(\exists x \in \mathbb{R}) [(\forall y \in \mathbb{R}) (xy \neq 1)].$$

• **Continuity of functions**: By the definition of continuity and using the logic symbol, *f* is continuous at a number *c* if

$$(\forall \varepsilon) (\exists \delta) \underbrace{(\forall x) [(|x - c| < \delta) \Rightarrow (|f(x) - f(c)| < \varepsilon)]}_{Q(\varepsilon, \delta)} .$$

- The universe for the variables ε and δ is the collection of positive real numbers. Therefore, sometimes we write
 (∀ ε > 0)(∃ δ > 0)(∀ x)[(|x c| < δ) ⇒ (|f(x) f(c)| < ε)].
- 2 The sentence $P(\varepsilon)$ is always true for any $\varepsilon > 0$.

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- The universe for the variables ε and δ is the collection of positive real numbers. Therefore, sometimes we write
 (∀ε > 0)(∃δ > 0)(∀x)[(|x c| < δ) ⇒ (|f(x) f(c)| < ε)].
- 2 The sentence $(\exists \, \delta) Q(\varepsilon, \delta)$ is always true for any $\varepsilon > 0$.

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• **Continuity of functions**: By the definition of continuity and using the logic symbol, *f* is continuous at a number *c* if

$$(\forall \varepsilon) (\exists \delta) \underbrace{(\forall x) [(|x-c| < \delta) \Rightarrow (|f(x) - f(c)| < \varepsilon)]}_{Q(\varepsilon, \delta)} .$$

- 2 The sentence $(\exists \delta)Q(\varepsilon, \delta)$ is always true for any $\varepsilon > 0$.
- Suppose ε is a given positive number. Then the truth set of Q(ε, δ) is non-empty which implies that "there is at least one positive number δ making the sentence Q(ε, δ) true".

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Definition

Two quantified statement are equivalent in a given universe if they have the same truth value in that universe. Two quantified sentences are equivalent if they are equivalent in every universe.

Example

Consider quantified sentences " $(\forall x)(x > 3)$ " and " $(\forall x)(x \ge 4)$ ".

- They are equivalent in the universe of integers because both are false.
- They are equivalent in the universe of natural numbers greater than 10 because both are true.
- O They are not equivalent in the universe X = [3.7,∞) of the real line.

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